

# Randomly Sampling from Orthonormal Matrices: Coherence and Leverage Scores

Ilse Ipsen

Joint work with Thomas Wentworth  
(thanks to Petros & Joel)

North Carolina State University  
Raleigh, NC, USA

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# This Talk

Given:

- Real  $m \times n$  matrix  $Q$  with **orthonormal columns**,  $Q^T Q = I_n$
- Real  $c \times m$  **“sampling” matrix**  $S$  with  $c \ll m$

Want: **Probability** that

- 1  $SQ$  has full column rank ( $\text{rank}(SQ) = n$ )
- 2 Condition number: Given  $\eta$

$$\kappa(SQ) = \|SQ\|_2 \|(SQ)^\dagger\|_2 \leq 1 + \eta$$

**Motivation: Blendenpik** [Avron, Maymounkov & Toledo 2010]

$\kappa(SQ)$  = condition number of preconditioned matrix

# Outline

- 1 Three different **uniform** sampling strategies
- 2 Bounds for **condition numbers** of sampled matrices  
*Based on **coherence***  
*Lower bounds on number of samples*
- 3 Improving on coherence: **Leverage scores**
- 4 Summary

# Three Different Uniform Sampling Strategies

Uniform sampling  $c$  rows out of  $m$

- 1 Sampling **without** replacement
- 2 Sampling **with** replacement (Exactly( $c$ ))
- 3 **Bernoulli** sampling

# Uniform Sampling without Replacement

[Gittens & Tropp 2011, Gross & Nesme 2010]

Choose random permutation  $k_1, \dots, k_m$  of  $1, \dots, m$

Sampling matrix  $S = \sqrt{\frac{m}{c}} \begin{pmatrix} e_{k_1}^T \\ \vdots \\ e_{k_c}^T \end{pmatrix}$

- $S$  is  $c \times m$ , and samples **exactly**  $c$  rows
- Each row sampled **at most once**
- Expected value  $\mathbf{E}(S^T S) = I_m$

# Uniform Sampling with Replacement (Exactly(c))

[Drineas, Kannan & Mahoney 2006]

**for**  $t = 1 : c$  **do**

    Sample  $k_t$  from  $\{1, \dots, m\}$  with probability  $1/m$   
    independently and **with replacement**

**end for**

Sampling matrix  $S = \sqrt{\frac{m}{c}} \begin{pmatrix} e_{k_1}^T \\ \vdots \\ e_{k_c}^T \end{pmatrix}$

- $S$  is  $c \times m$ , and samples **exactly**  $c$  rows
- A row can be sampled **more than once**
- Expected value  $\mathbf{E}(S^T S) = I_m$

# Bernoulli Sampling

[Avron, Maymounkov & Toledo 2010, Gittens & Tropp 2011]

$$S = 0_{m \times m}$$

**for**  $j = 1 : m$  **do**

$$S_{jj} = \sqrt{\frac{m}{c}} \begin{cases} 1 & \text{with probability } \frac{c}{m} \\ 0 & \text{with probability } 1 - \frac{c}{m} \end{cases}$$

**end for**

- $S$  is  $m \times m$ , and samples each row **at most once**
- **Expected** number of sampled (non zero) rows:  $c$
- Expected value  $\mathbf{E}(S^T S) = I_m$

# Sampling Rows from Orthonormal Matrices

Sampling  $c$  rows from  $m \times n$  matrix  $Q$  with  $Q^T Q = I_n$   
 $m = 10^4$ ,  $n = 5$  (30 runs for each value of  $c$ )

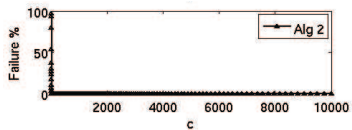
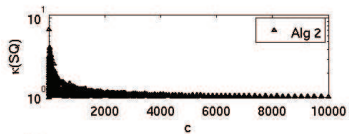
For each sampling strategy:

- 1 Two-norm condition number of  $SQ$   
 $\kappa(SQ) = \|SQ\|_2 \|(SQ)^\dagger\|_2$  (if  $SQ$  has full column rank)
- 2 Percentage of matrices  $SQ$  that are rank deficient

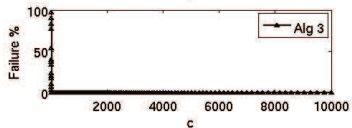
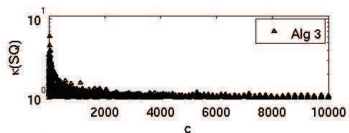


# First Comparison

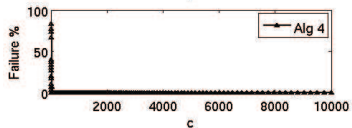
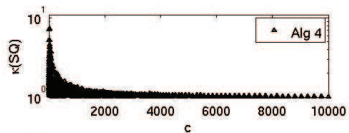
Sampling without replacement



Sampling with replacement (Exactly(c))

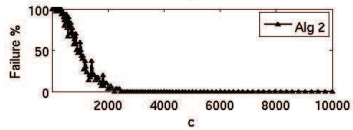
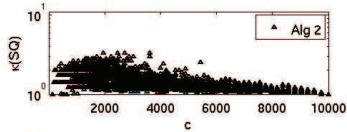


Bernoulli sampling

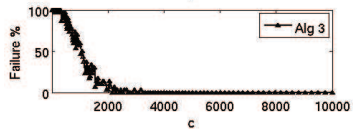
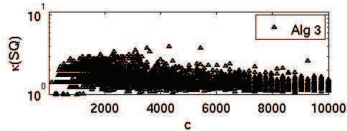


# Second Comparison

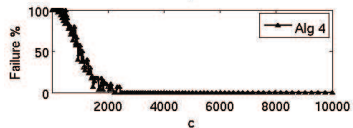
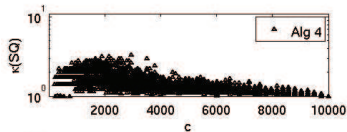
Sampling without replacement



Sampling with replacement (Exactly( $c$ ))



Bernoulli sampling



# Comparison of Sampling Strategies

Three strategies for sampling  $SQ$ :

*Sampling without replacement*

*Sampling with replacement (Exactly( $c$ ))*

*Bernoulli sampling*

## Summary

*Little difference among the sampling strategies*

*If  $SQ$  has full rank then  $\kappa(SQ) \leq 10$*

**Recommendation:** Use sampling **with** replacement

*Fast: Need to generate/inspect only  $c$  values*

*Easy to implement*

*Replacement does not affect accuracy*

*(for small amounts of sampling)*

# Condition Number Bound

Same bound for all three sampling strategies

- $m \times n$  matrix  $Q$  with  $Q^T Q = I_n$   
coherence  $\mu \equiv \max_{1 \leq j \leq m} \|e_j^T Q\|_2^2$
- Number of samples  $c \geq n$ ,  $0 < \epsilon < 1$
- Failure probability

$$\delta \equiv n \left( f(-\epsilon)^{c/(m\mu)} + f(\epsilon)^{c/(m\mu)} \right)$$

where  $f(x) = e^x(1+x)^{-(1+x)}$

With probability at least  $1 - \delta$

$$\kappa(SQ) \leq \sqrt{\frac{1+\epsilon}{1-\epsilon}}$$

# Proofs

- 1 Sampling **without** replacement: [Tropp: *SRHT*, 2011]
- 2 Sampling **with** replacement:  
Use matrix Chernoff bound from [Tropp: *Userfriendly*, 2011]
- 3 **Bernoulli** sampling:  
Use matrix Chernoff bound from [Tropp: *Userfriendly*, 2011]  
See also [Gittens & Tropp, 2011]

# Lower Bounds on Number of Samples

Failure probability

$$\delta \leq n \left( f(-\epsilon)^{c/(m\mu)} + f(\epsilon)^{c/(m\mu)} \right)$$

where  $f(x) = e^x(1+x)^{-(1+x)}$

Number of samples

$$c \geq \mu m \ln(2n/\delta) / |\ln(f(\epsilon))|$$

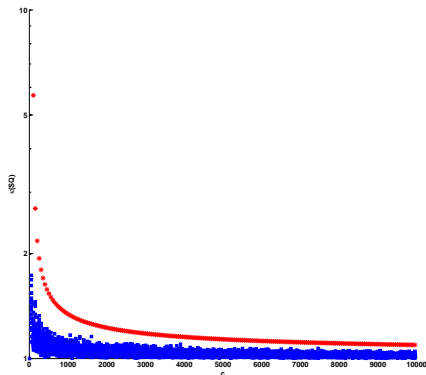
- Condition number: For  $\kappa(SQ) \leq 10$  need  $\epsilon = 99/101$

$$c \geq 2.7 \mu m \ln(2n/\delta)$$

- If  $Q$  has minimal coherence  $\mu = n/m$

$$c \geq 2.7 n \ln(2n/\delta)$$

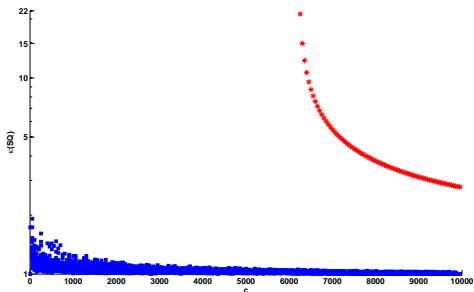
# Tightness of Bound: Low Coherence ( $\mu = 1.5n/m$ )



Numerical bound for  $c \geq 104$ , lower bound  $c \geq 108$

( $m = 10^4$ ,  $n = 4$ ,  $\delta = .01$ )

# Tightness of Bound: Higher Coherence ( $\mu = 100n/m$ )



Numerical bound  $c \geq 6,254$ , lower bound  $c \geq 7,219$   
( $m = 10^4$ ,  $n = 4$ ,  $\delta = .01$ )



# Conclusions from the Bound

Input:  $m \times n$  matrix  $Q$  with orthonormal columns

- Bound is based on coherence  $\mu = \max_{1 \leq j \leq m} \|e_j^T Q\|_2^2$
- Predicts correct magnitude for condition number of sampled matrix, even for small matrix dimensions
- Required number of samples  $c \geq m \mu \ln n$   
But too pessimistic for matrices with higher coherence
- Informative only for matrices  
*that are very tall and skinny,  $m \gg n$   
with almost minimal coherence  $\mu \approx n/m$*

# Improving on Coherence: Leverage Scores

Idea: Use **all** row norms

- $Q$  is  $m \times n$  with orthonormal columns
- **Leverage scores** = row norms<sup>2</sup>

$$\ell_j = \|e_j^T Q\|_2^2, \quad 1 \leq j \leq m$$

[Hoaglin & Welsch 1978], [Chatterjee & Hadi 1986]

[Drineas, Mahoney & al], [Avron & Toledo]

- **Coherence**  $\mu = \max_j \ell_j$
- **Low coherence**  $\approx$  uniform leverage scores

# Leverage Score Bound for Sampling with Replacement (Exactly(c))

- $m \times n$  matrix  $Q$  with orthonormal columns
- Leverage scores  $\ell_j = \|e_j^T Q\|_2^2$ ,  $\mu = \max_{1 \leq j \leq m} \ell_j$

$$L = \text{diag}(\ell_1 \quad \dots \quad \ell_m)$$

- $0 < \epsilon < 1$

Failure probability

$$\delta = 2n \exp\left(-\frac{3}{2} \frac{c \epsilon^2}{m(3 \|Q^T L Q\|_2 + \mu \epsilon)}\right)$$

With probability at least  $1 - \delta$ :  $\kappa(SQ) \leq \sqrt{\frac{1+\epsilon}{1-\epsilon}}$

## Leverage Scores vs. Coherence

- Failure probability

$$\delta = 2n \exp\left(-\frac{3}{2} \frac{c \epsilon^2}{m (3 \|Q^T L Q\|_2 + \mu \epsilon)}\right)$$

- Bounds in terms of coherence:

$$\mu^2 \leq \|Q^T L Q\|_2 \leq \mu$$

- Estimation in terms of largest leverage scores

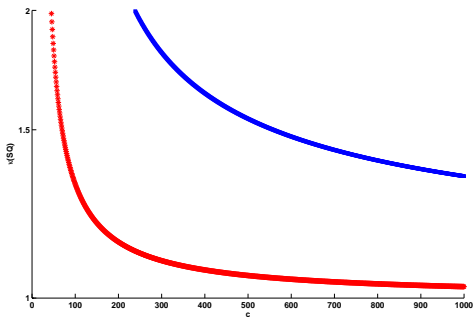
If  $k = 1/\mu$  is an integer then

$$\|Q^T L Q\|_2 \leq \mu \sum_{j=1}^k \ell_{[j]}$$

where  $\ell_{[1]} \geq \dots \geq \ell_{[m]}$

# Improvement with Leverage Score Bound

Low coherence:  $\mu = 1.5n/m$ , small amounts of sampling

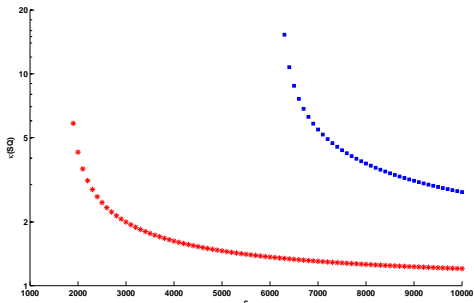


Leverage score bound vs. Coherence bound

( $m = 10^4$ ,  $n = 4$ ,  $\delta = .01$ )

# Improvement with Leverage Score Bound

Higher coherence:  $\mu = 100n/m$ , large amounts of sampling

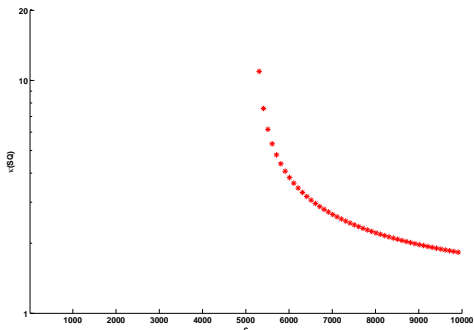


Leverage score bound vs. Coherence bound

( $m = 10^4$ ,  $n = 4$ ,  $\delta = .01$ )

# Improvement with Leverage Score Bound

More columns, higher coherence:  $\mu = 100n/m$



Coherence bound does not hold  
( $m = 10^4$ ,  $n = 10$ ,  $\delta = .01$ )

# Summary

- Sampling rows from  $m \times n$  matrices with orthonormal columns
- Want: Bounds for **condition number** of sampled matrices
- **Three different sampling strategies:**  
Without replacement, with replacement, Bernoulli
- Numerical experiments:  
All strategies behave the same for  $c < m$
- **Recommendation:** Sampling **with** replacement
  - Fast: Need to generate/inspect only  $c$  values*
  - Easy to implement*
  - Replacement does not affect accuracy*  
*(for small amounts of sampling)*



## Summary (cdt)

- Bounds for condition number of sampled matrices

*Explicit and non-asymptotic*

*Predictive even for **small matrix dimensions***

- Bound based on coherence  $\mu$ :

*Same bound for all strategies*

*Number of samples  $c \geq \mu m \ln n$*

- Bound based on leverage scores:

*Tighter than coherence-based bound*

*Holds for **smaller number of samples  $c$** ,*

*higher coherence and fatter matrices*